1	TITLE OF THE INVENTION
2	QoS-based Shortest Path Routing for a Hierarchical Communication
3	Network
4	BACKGROUND OF THE INVENTION
5	Field of the Invention
6	The present invention relates generally to communications networks,
7	and more specifically to an on-demand QoS (quality-of-service)-based
8	routing for a hierarchical communication network.
9	Description of the Related Art
lO	RFC (Request for Comments) 2328 and 2676 texts describe a
1	hierarchical communication network in which QoS-based on-demand routing
12	is performed using the OSPF (Open Shortest Path First) algorithm, known as
13	QOSPF (QoS extended OSPF). On-demand QoS routing algorithm is one that
4	determines a QoS route to a user-specified destination using the Dijkstra
15	algorithm. This QoS routing is particularly useful for QoS-guaranteed
6	networks such as multi-protocol label switching (MPLS) networks. The
17	hierarchical communication network is comprised of a plurality of routers
8	interconnected by links. Each of the routers belongs to one of a plurality of
19	areas, one of which is the backbone area which is traversed by traffic between
20	non-adjacent areas. Adjacent areas are interconnected by at least one router
21	known as an area border router (ABR). If an area border router receives an
22	on-demand QoS route calculation request from a user, requesting a route to a
23	destination that is located in one of its neighboring areas, the router calculates
24	a QoS-based shortest path tree (SPT) according to the Dijkstra algorithm.
25	However since the router has no knowledge of which areas can be traversed

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1	to reach the specified destination, the QoS-SPT calculation must be
2	performed for all of its neighboring areas. Further, the router has no
3	knowledge of which remote area border routers can be used as intermediate
4	routers to reach a remote destination. Therefore, if the destination is in a
5	remote area and can be reached via the backbone area, the QoS-SPT must be
6	calculated for all possible routes of the backbone area from the source router
7	to the remote area border routers, in addition to the QoS-SPT calculations for
8	all possible routes of the local area of the source router. As a result, the prior
9	art routing technique is wasteful of QoS-SPT calculations.
10	SUMMARY OF THE INVENTION
11	It is therefore an object of the present invention to provide a
12	hierarchical communication network and a method of communication that
13	eliminate the wasteful route calculations.
14	Another object of the present invention is to provide QoS-based
15	routing that allows each router of the network to possess the knowledge of all
16	areas that can be traversed in advance of selecting links.
17	A further object of the present invention is to provide QoS-based
18	routing that allows each router of the network to possess, for each traversable
19	area, the knowledge of all area border routers that can be used as transit
20	routers to reach a remote destination via the traversable area.
21	According to a first aspect of the present invention, there is provided a
22	router for a hierarchical communication network which is divided into a
23	plurality of areas in each of which a plurality of the router are interconnected
24	by links, comprising a first table having a plurality of entries respectively

corresponding to reachable destinations, each of the entries including an

- 1 intra-area or an inter-area indication and an area identifier identifying at least
- 2 one traversable area and a plurality of second tables respectively
- 3 corresponding to the areas, each of the second tables holding quality-of-
- 4 service (QoS) values of the links of the corresponding area. A processor is
- 5 responsive to a request signal specifying a destination and a QoS value for
- 6 making reference to one of the entries of the first table and one of the second
- 7 tables corresponding to the specified destination, selecting links of the area
- 8 identified by the area identifier of the referenced entry which links satisfy the
- 9 specified QoS value, and performing a calculation according to a shortest
- 10 path finding algorithm on the selected links to find a shortest path to the
- 11 specified destination if the intra-area indication is included in the referenced
- 12 entry, or performing the shortest path calculation on the selected links to find
- a shortest path tree in the identified area and determining a route from the
- 14 shortest path tree.
- 15 According to a second aspect, the present invention provides a router
- 16 for a hierarchical communication network which is divided into a plurality of
- 17 areas in each of which a plurality of the router are interconnected by links,
- 18 wherein neighboring ones of the areas are interconnected by at least one area
- 19 border router. The router comprises a first table having a plurality of entries
- 20 respectively corresponding to reachable destinations, each of the entries
- 21 including an intra-area or an inter-area indication, an area identifier
- 22 identifying at least one traversable area, and a list of area border routers if the
- 23 inter-area indication is included and a plurality of second tables respectively
- 24 corresponding to the areas, each of the second tables holding quality-of-
- 25 service (QoS) values of the links of the corresponding area. A processor is

value.

1	responsive to a request signal specifying a destination and a QoS value for
2	making reference to one of the entries of the first table and one of the second
3	tables corresponding to the specified destination, selecting links of the area
4	identified by the area identifier of the referenced entry which links satisfy the
5	specified QoS value, and performing a calculation according to a shortest
6	path finding algorithm on the selected links to find a shortest path to the
7	specified destination if the intra-area indication is included in the referenced
8	entry, or performing the shortest path calculation on the selected links until a
9	shortest path tree is found for all routers of the list of the referenced entry or
10	until an end of the calculation is reached when the tree is not found for all the
11	routers if the inter-area indication is included in the referenced entry, and
12	determining from the shortest path tree a route having an optimum QoS

According to a third aspect of the present invention, there is provided a hierarchical communication network which is divided into a plurality of areas in each of which a plurality of the router are interconnected by links. Each of the routers comprises a first table having a plurality of entries respectively corresponding to reachable destinations, each of the entries including an intra-area or an inter-area indication and an area identifier identifying at least one traversable area and a plurality of second tables respectively corresponding to the areas, each of the second tables holding quality-of-service (QoS) values of the links of the corresponding area. A processor is responsive to a request signal specifying a destination and a QoS value for making reference to one of the entries of the first table and one of the second tables corresponding to the specified destination, selecting links of

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1	the area identified b	the area identifier of the referenced ent	ry which links
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- 2 satisfy the specified QoS value, and performing a calculation according to a
- 3 shortest path finding algorithm on the selected links to find a shortest path to
- 4 the specified destination if the intra-area indication is included in the
- 5 referenced entry, or performing the shortest path calculation on the selected
- 6 links to find a shortest path tree in the identified area and determining a
- 7 route from the shortest path tree.

According to a fourth aspect, the present invention provides a hierarchical communication network which is divided into a plurality of areas in each of which a plurality of routers are interconnected by links, wherein neighboring ones of the areas are interconnected by at least one area border router. Each of the routers comprises a first table having a plurality of entries respectively corresponding to reachable destinations, each of the entries including an intra-area or an inter-area indication, an area identifier identifying at least one traversable area, and a list of area border routers if the inter-area indication is included, and a plurality of second tables respectively corresponding to the areas, each of the second tables holding quality-ofservice (QoS) values of the links of the corresponding area. A processor is responsive to a request signal specifying a destination and a QoS value for making reference to one of the entries of the first table and one of the second tables corresponding to the specified destination, selecting links of the area identified by the area identifier of the referenced entry which links satisfy the specified QoS value, and performing a calculation according to a shortest path finding algorithm on the selected links to find a shortest path to the specified destination if the intra-area indication is included in the referenced

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entry, or performing the shortest path calculation on the selected links until a 1 shortest path tree is found for all routers of the list of the referenced entry or 2 until an end of the calculation is reached when the tree is not found for all the 3 routers if the inter-area indication is included in the referenced entry, and 4 determining from the shortest path tree a route having an optimum QoS 5 value. б According to a fifth aspect of the present invention, there is provided a 7 routing method for a hierarchical communication network which is divided 8 into a plurality of areas in each of which a plurality of the router are 9 interconnected by links, each of the routers comprising a first table having a 10 plurality of entries respectively corresponding to reachable destinations, each 11 of the entries including an intra-area or an inter-area indication and an area 12 identifier identifying at least one traversable area, and a plurality of second 13 tables respectively corresponding to the areas, each of the second tables 14 holding quality-of-service (QoS) values of the links of the corresponding area, 15 each of the routers functioning as a source router when a request signal is 16 received. The method comprises the steps of receiving, at the source router, a 17 request signal specifying a destination and a QoS value and making reference 18 to one of the entries of the first table and one of the second tables 19 corresponding to the specified destination, selecting links of the area 20 identified by the area identifier of the referenced entry which links satisfy the 21 specified QoS value, and performing a calculation according to a shortest 22 path finding algorithm on the selected links to find a shortest path to the 23 specified destination if the intra-area indication is included in the referenced 24

entry, or performing the shortest path calculation on the selected links to find

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a shortest path tree in the identified area and determining a route from the shortest path tree.

According to a sixth aspect, the present invention provides a routing method for a hierarchical communication network which is divided into a plurality of areas in each of which a plurality of routers are interconnected by links, the routers of neighboring areas being interconnected by at least one area border router, wherein each of the routers functions as a source router when a request signal is received and includes a first table having a plurality of entries respectively corresponding to reachable destinations, each of the entries including an intra-area or an inter-area indication, an area identifier identifying at least one traversable area, and a list of area border routers if the inter-area indication is included, and a plurality of second tables respectively corresponding to the areas, each of the second tables holding quality-ofservice (QoS) values of the links of the corresponding area. The routing method comprises the steps of receiving, at the source router, a request signal specifying a destination and a QoS value, for making reference to one of the entries of the first table and one of the second tables corresponding to the specified destination, selecting links of the area identified by the area identifier of the referenced entry which links satisfy the specified QoS value, and performing a calculation according to a shortest path finding algorithm on the selected links to find a shortest path to the specified destination if the intra-area indication is included in the referenced entry, or performing the shortest path calculation on the selected links until a shortest path tree is found for all routers of the list of the referenced entry or until an end of the calculation is reached when the tree is not found for all the routers if the

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1	inter-area indication is included in the referenced entry, and determining
2	from the shortest path tree a route having an optimum QoS value.
3	Due to the listing of the area ID in the first table, the path finding
4	calculation for intra-area destinations is limited only to the local area.
5	Wasteful calculations on unnecessary links for other areas are eliminated.
6	Further, due to the listing of at least one traversable area ID and the
7	router ID's of corresponding area border routers in the first table, the path
8	finding calculation for inter-area destinations is limited only to the
9	traversable area. Wasteful calculations on unnecessary links for other areas
10	are eliminated. In addition, the amount of shortest path tree calculations is
11	minimized due to the fact that the calculation is performed until a QoS
12	shortest path tree is found for all area border routers of the traversable area or
13	until it terminates of its own accord when such a path is not found for all area
14	border routers.
15	BRIEF DESCRIPTION OF THE DRAWIGNS
16	The present invention will be described in detail further with reference
17	to the following drawings, in which:
18	Fig. 1 is a block diagram of an exemplary hierarchical communication
19	network of the present invention;
20	Fig. 2 is a block diagram of a representative router of the network of
21	Fig. 1;
22	Figs. 3A, 3B and 3C are illustrations of resource tables of the
23	representative router; and
24	Figs. 4A and 4B are flowcharts of the operation of the processor of Fig.

2 according to the present invention.

1	DETAILED DESCRIPTION
2	In Fig. 1 an IP (Internet Protocol) network is illustrated in simplified
3	form in which routing is performed based on a user-requested QoS (quality
4	of service) value in accordance with the present invention. As a typical
5	example , the QOSPF (QoS extended Open Shortest Path First) algorithm will
6	be explained. The IP network 1 is comprised of a hierarchical QOSPF
7	network 2 and an autonomous system 3. The QOSPF network 2, which is
8	also an autonomous system in the IP network, is formed by three OSPF areas
9	4, 5 and 6, with the area 5 being a backbone area that functions as a core of
10	the QOSPF network. In each of these areas, routers are interconnected as
11	neighbors sharing the same Area identifier (ID) and the routers on the border
12	of two adjacent areas operate as Area Border Routers (ABR). As illustrated,
13	the area 4 is comprised of routers 41 to 45, with the routers 44 and 45 being
14	ABRs connected to routers 51, 52 of the backbone area 5 and the router 41
15	being connected to a network 40 as a neighbor of the area 4. Area 6 is
16	comprised of routers 61 to 65, with the routers 61 and 62 being ABRs
17	connected to routers 52, 53 of the backbone area 5 and the router 65 being
18	connected to a network 60 as a neighbor of the area 6. Routers 42 and 51 are
19	autonomous system border routers (ASBRs) for the autonomous system 3.
20	Within the OSPF network, the routers send routing updates with the use of
21	link-state advertisement packets, or LSAs such as router LSA, network LSA,
22	summary LSA and AS-external LSA.
23	Each of the area border routers 44 and 45 shrinks the routing updates
24	of their local area 4 into a summary and sends it to the backbone 5 and
25	shrinks the routing updates of backbone area 5 into a summary for

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distribution within their local area 4. Likewise, each of the area border 1 routers 61 and 62 shrinks the routing updates of their local area 6 into a 2 summary for distribution to the backbone area 5 and shrinks the routing 3 updates of backbone 5 into a summary for distribution within their local area 4 6. Note that the summary of backbone 5 distributed within the area 4 also 5 contains the summary of area 6. Hence all routers of area 4 have the 6 knowledge of which destinations are reachable within area 6 as well as 7 within the backbone area 5. Likewise, all routers of area 6 have the 8 knowledge of which destinations are reachable within areas 4 as well as 9 within the backbone area 6. 10 As shown in Fig. 2, each of the routers of the present invention 11 includes an interface 20 connected via communication links to neighboring 12 routers. The interface 20 performs routing with the neighbors according to 13 the routing protocol of the OSPF domain. Interface 20 is associated with a 14 topology table 21 and a plurality of resource tables to maintain network 15 database by exchanging LSAs with neighboring routers. As a representative 16 router, the router 44 may includes a resource table 22 for holding the 17 18 bandwidth database of its local area 4, a resource table 23 for holding the bandwidth database of the backbone area 5 and a resource table 24 for 19 holding a summarized database of the non-adjacent area 6. More specifically, 20 the summarized resource table 24 contains hot count values and remaining 21 bandwidths of routes from the area border routers 61 and 62 to the network 22 23 60.

Topology table 21 has a number of entries respectively corresponding

to a plurality of reachable destinations. Each entry is subdivided into a

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in the ABR list.

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plurality of fields including an IN/OUT field, an AREA ID field, and an ABR 1 LIST field. The IN/OUT field indicates whether the destination of the entry 2 is inside or outside of the local area of the router. The AREA ID field contains 3 the Area IDs of all areas that can be traversed along routes to the destination. 4 The ABR LIST field indicates one or more area border routers (ABRs) along 5 possible routes to the reachable destination. б 7 Router 44, for example, uses LSA packets to create entries for the networks 40 and 60 and the autonomous system 3 in the topology table 21. 8 9 Specifically, the router 44 examines the router LSA and the network LSA flooded in the local area 4 and recognizes that the network 40 exists within 10 11 the same area 4 as router 44 and sets an "IN" (intra-area) indication in the 12 IN/OUT field of the first entry of the topology table 21 and sets ID = 4 in the 13 AREA ID field and leaves the ABR LIST field of this entry vacant. Router 44 examines the router LSAs and network LSAs flooded in the areas 4 and 5 and 14 determines that the network 60 is not the same member of the local area 4 15 16 and proceeds to examine the summary LSA advertised to the backbone 5 17 from the routers 61 and 62 and sets an "OUT" (inter-area) indication and ID = 5 in the IN/OUT and AREA ID fields of the second entry and sets Routers 61 18 19 and 62 in the ABR LIST field. In the case of the autonomous system 3, the 20 router 44 determines that it is not the same member of the area 4 from the 21 router LSA and network LSA flooded in the local area 4 and the backbone 22 area 5 and proceeds to refer to the AS-external LSAs advertised to the OSPF 23 network 2 and sets an "OUT" indication in the IN/OUT field of the third

entry, and ID = 4 and ID = 5 in the AREA ID field, and sets Routers 42 and 51

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Figs. 3A, 3B and 3C show details of the resource tables 22, 23 and 24. 1 To create these resource tables, the router 44 uses resource data of active links 2 of area 4 and stores their usable bandwidths for both outgoing and incoming 3 links in the resource table 22. In the same manner, the router 44 stores usable 4 bandwidths of active links of backbone area 5 in the resource table 23. Router 5 44 is advertised of summarized resource data of active links in the area 6 6 from the routers 61 and 62 as shown in Fig. 3C. 7 A processor 25, is connected to the tables 20, 21, 22 and 23. As will be 8 described, the processor 25 is responsive to a request from users to perform 9 on-demand QoS route calculations using the contents of the topology and 10 resource tables and replies with a return message containing the result of the 11 12 route calculations. The operation of the processor 25 proceeds according to the flowcharts 13 14 shown in Figs. 4A and 4B. When an OSPF router receives a request signal from a user for an on-15 demand QoS route calculation (step 70), the processor 25 proceeds to step 71 16 to make reference to an entry of the topology table 21 that corresponds to a 17 18 destination specified in the request signal and reads the IN/OUT field to 19 determine whether the specified destination is inside or outside of the local area of the source router. 20 At step 72, the processor reads the Area ID of the referenced entry of 21 the topology table. At step 73, the processor makes reference to one of the 22 resource tables that corresponds to the area identified by the read Area ID 23 and reads resource and routing data. At step 74, the processor 25 uses the 24 read routing data to select links whose bandwidths satisfy a value specified 25

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in the request signal.

At step 75, the processor 25 performs calculations according to the
Dijkstra algorithm on the selected links to find a QoS shortest path to the
specified destination. If a shortest path is not found (step 76), a reply
message is sent to the requesting user to inform that the destination is
unreachable. If the decision is affirmative at step 76, the processor proceeds
to step 77 to inform the user of the routing information of the calculated
shortest path.

If the destination specified in the request signal is outside of the local area, flow proceeds from step 71 to step 80 (Fig. 4B) to read the Area ID of the referenced entry of the topology table. Processor 25 then makes reference to one of the resource tables that corresponds to the read Area ID (step 81), and uses the read routing data to select links whose bandwidths satisfy the user-specified value (step 82). At step 83, the processor reads all Router ID's of the ABR list of the referenced entry of the topology table.

At step 84, the processor 25 performs calculations according to the Dijkstra algorithm on the selected links to find a QoS shortest path tree to all routers of the ABR list. If a shortest path tree is found for all routers of the ABR list (step 85), the processor terminates the calculation at step 86 and proceeds to step 87. If a shortest path tree is not found for all routers the ABR list, the processor proceeds from step 85 to step 88 to check to see if the calculation has terminated. If not, flow returns to step 84 to continue the calculation. Therefore, if the decision at step 88 is affirmative it can be determined that a shortest path tree has not been found for any router of the ABR list or one has been found for some of the routers of the ABR list.

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At step 87, the processor determines whether the above process has been performed on all traversable areas identified by the Area ID's of the referenced entry of the topology table. If not, flow returns to step 80 to read the next Area ID from the referenced entry and repeats until the shortest path finding calculation is performed on QoS-satisfying links of all traversable areas indicated in the referenced entry of the topology table. Decision step 89 determines whether at least one shortest path tree has been found. If the decision is affirmative, the processor selects a route with a maximum remaining bandwidth from the shortest path tree (step 90) and informs the requesting user of the selected route (step 91). Otherwise, the processor sends a message indicating that the destination is unreachable. Assume that the processor 25 of router 44 receives an on-demand QoS route calculation request from a user, requesting a 10-Mbps route and specifying the network 40 as a destination (step 70). Processor 25 first looks up the topology table 21 and determines that the destination is in the same local area (step 71). In the topology table 21, the entry of network 40 is referenced for reading the intra-area indication and the Area ID = 4 (step 72) and the resource table 22 is referenced corresponding to the Area ID = 4 (step 73) for reading the resource and routing data of the local area 4. Links of the area 4 whose remaining bandwidths satisfy the requested 10 Mbps are selected (step 74). Thus, the 5-Mbps outgoing link from router 42 to router 41 is excluded in the link selection process and the processor performs Dijkstra algorithm path finding calculation (step 75) on the selected links to find a route 100 as a shortest path to the destination (see Fig. 1), including the first link from router 44 to router 42, the intermediate link from router 42 to router

 1 43 and the final link from router 43 to router 41.

Thus, due to the listing of the area ID in the topology table, the path

3 finding calculation for intra-area destinations is limited only to the local area.

4 Wasteful calculations on unnecessary links for other areas are eliminated.

5 If the user requests a 15-Mbps route to the network 60 (step 70).

6 Processor 25 determines that the destination is outside of the local area (step

7 71). Processor 25 then examines the Area-ID field of the entry and knows that

the backbone area 5 is the traversable area and the network 60 can be reached

9 via the backbone area 5. In the topology table, the entry of network 60 is

referenced and the Area ID = 5 of the backbone area 5 is read (step 80).

11 Processor 25 knows that the network 60 can be reached via links of the

backbone area 5 to the routers 61 and 62. corresponding to Area ID = 5, the

13 resource table 23 is referenced (step 81) and links of remaining bandwidth of

14 at least 15-Mbps are selected from this resource table (step 82). Processor 25

reads the router identifiers ID = 61 and ID = 62 of the ABR list of the

16 referenced entry (step 83). Processor 25 performs the Dijkstra algorithm

17 calculation on the selected links to find a shortest path tree that extends from

the source router 44 to the area border routers 61 and 62 (steps 84 to 87). For

example, two routes 101 and 102 from the router 44 to area border routers 61

20 and 62 are selected. Route 101 includes a first link from router 44 to router

21 51, an intermediate link from router 51 to router 53, an intermediate link from

22 router 53 to router 52 and a final link from router 52 to router 61. Second

23 route 102 includes a first link from router 44 to router 51, an intermediate link

24 from router 51 to router 53 and a final link from router 53 to router 62. Since

25 there is only one Area ID in the entry of the network 60, the processor then

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proceeds from step 87 to step 89. Since two routes are determined, the 1 processor examines the summarized resource table 24 and compares the two 2 selected routes in terms of bandwidth available to the network 60 in the area 3 6 or hop count values of routes from the source router 44 to the area border 4 routers 61 and 62. 5 Thus, due to the listing of at least one traversable area ID and the 6 router ID's of corresponding area border routers in the topology table, the 7 path finding calculation for inter-area destinations is limited only to the 8 traversable area. Wasteful calculations on unnecessary links for other areas 9 are eliminated. Further, the amount of shortest path tree calculations is 10 minimized due to the fact that the calculation is performed until a shortest 11 path tree is found for all area border routers of the traversable area or until it 12 terminates of its own accord when such a path is not found for all area border 13 14 routers. If the policy of the OSPF network places priority on bandwidth, the 15 processor makes a decision in favor of the route from the router 61 to the 16 destination because of its greater remaining bandwidth than the route from 17 router 62 to the same destination. Therefore, the requesting user is informed 18 of the route 101 as a best route. If the routes 101 and 102 have different 19 values of minimum bandwidth, the larger of these will also be taken into 20 21 account in the final process of route selection along with the bandwidths 22 available in the area 6. If the policy of the OSPF network places priority on hop count value, 23 the processor produces a first sum of the hop count of route 101 plus the hop 24

count of the route from router 61 to the destination and a second sum of the

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- 1 hop count of route 102 plus the hop count of the route from router 62 to the
- destination. Since the first sum equals 7 (= 4 + 3) and the second sum equals
- 6 (= 3 + 3), the processor makes a decision in favor of route 102 because of its
- 4 smaller total value of hop count to the network 60.